

Comp 310 Computer Systems and Organization

Lecture #20 File Systems (Implementation – Part 2)

Prof. Joseph Vybihal

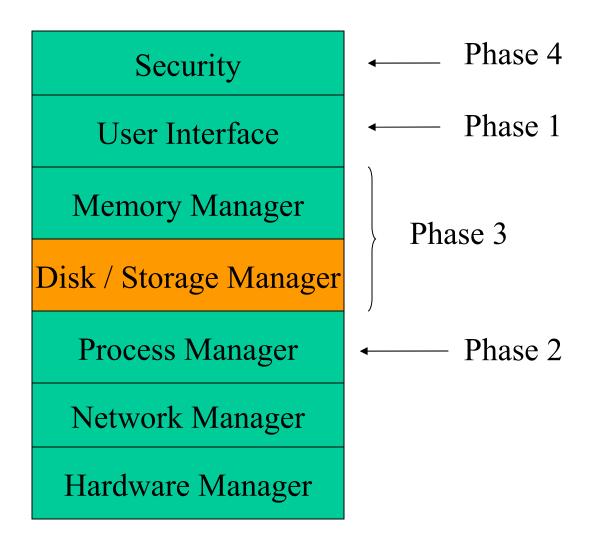


### Announcements

- Kelvin teaching next class
  - Nov 18



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Part 1

### File System Mounting



# Mounting (a.k.a Mapping)

- We cannot simply assume a directory is available to us
- Some devices like diskettes and hard disks have their directories available by default
- Other directory trees are not:
  - Private folders
  - Network folders
- To get access to them we need to perform an operation called Mounting or Mapping



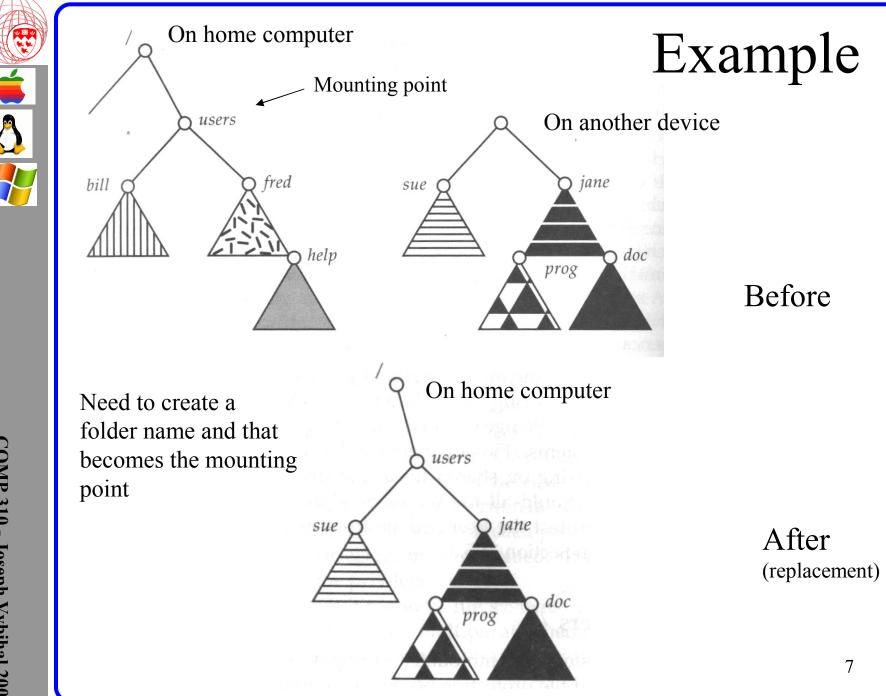
## MSDOS Example

c:\>

```
d:billsmith/>
```

```
c: \geq map d:
```

```
\leftarrow un-mount
```



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 mkdir /media/floppy0  $\leftarrow$  a mount point \$ mkfs deviceName options mkfs -t msdos /dev/fd0  $\leftarrow$  making a file system \$ mount existingDeviceName existingDirectory mount /dev/fd0 /media/floppy0  $\$  cat /etc/mnttab  $\leftarrow$  see the mounted directories \$ unmount f10

Unix Example



# How does mounting work?

- Two concepts
  - Directory Trees and FATs
  - Getting directories



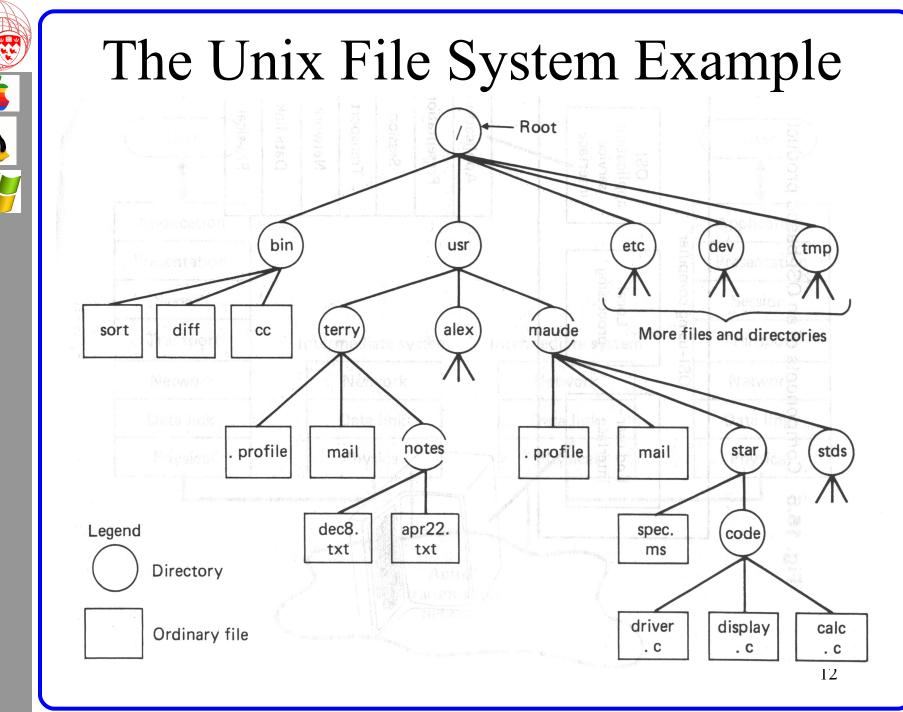
# FATs vs Directory Trees

- FAT
  - A structure on disk that records the folders and files on disk.
  - A semi-static data structure
- Directory Tree
  - A data structure in RAM (table or linked)
  - Generated at login time
    - System tree (can be hidden)
    - User tree (can limit the view of the hard disk)

# Getting Directories

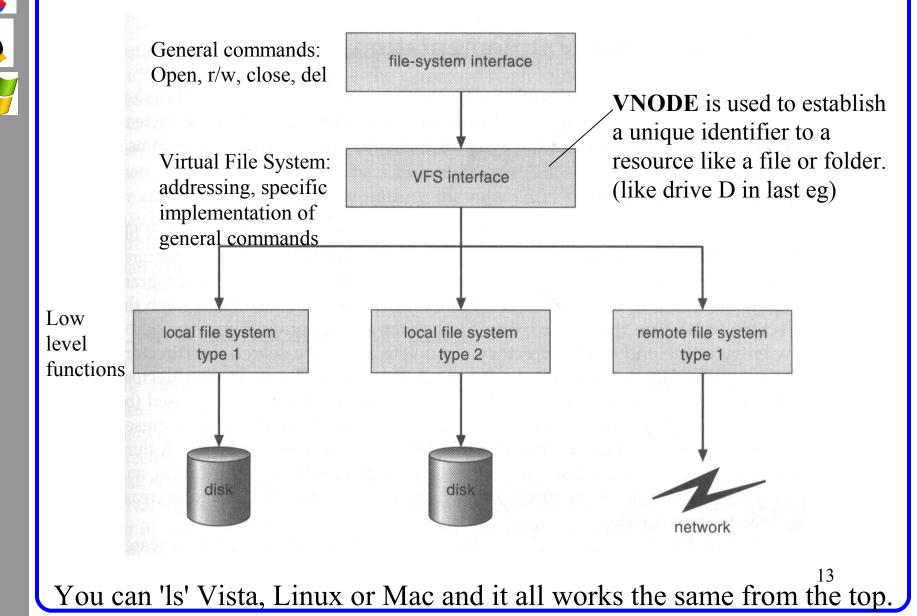
- If the directory tree is implemented as a linked list (tree)
  - Mounting is adding a node to the linked list
  - It then appears as an entry when 'ls' and 'cd' used
- If the directory tree is implemented through the FAT then the mounting point requires a new row entry, and that would look like a file.

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# Virtual File Systems



# Remote Mounting Methods

- Distributed File System (DFS)
  - Remote directories are made visible though software
  - Must login with user name and password
- Network File System (NFS)
  - Network directory mounted on a drive letter after login validation and verification that user's ID is logged as a valid member/user of the directory.
    - Local directory x  $\rightarrow$  network sever address : / usr
- Domain Name System (DNS)
  - Host machine name mapped to internet address (local):
    - Local directory  $x \rightarrow mimi.cs.mcgill.ca$  : /usr
  - Then it functions much like NFS
- Network Information Service (NIS) ~Sun
  - Unsecured centralize database of all servers, users, printers and directories.
  - Uses internet IP addresses like DNS

# File Security Methods

- File system records the following for every file and folder:
  - Owner,
  - Read, Write, Execute privileges by:
    - Owner
    - Friends (general or can set groups)
    - Public
  - Mountable if you have access right
    - User name and password on server of directory
    - You have r/w, or x privilege



## Part 2

#### File System Implementation

# **RAM Directory Implementations**

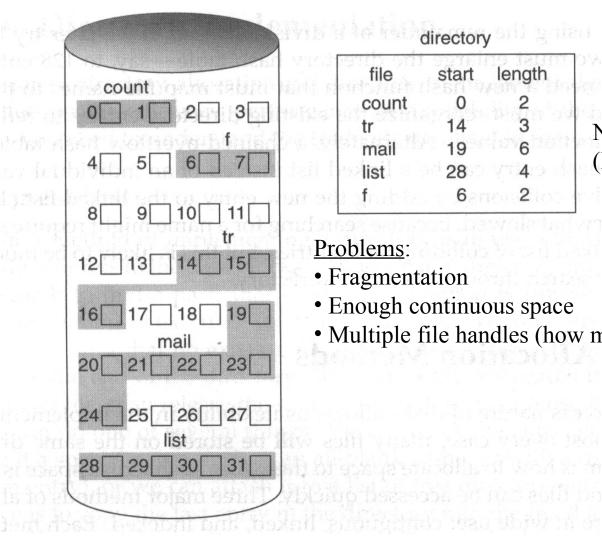
- P Tabular
  - A linear list of available files
  - One table for each folder
  - Each entry is a file or pointer to another table
  - Locate files using a linear search O(n) search, O(1) insert
- Hash Tables
  - Similar to tabular but files are located using a hash function to index to the file
  - Hash collisions are handled by linear searches from the point of the crash down the list. O(1) or O(n)
- B-Tree
  - Directory stored as a Tree
  - $O(\log n)$  search,  $O(n \log n)$  insertion

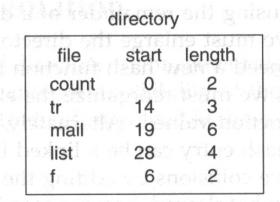


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### File Allocation Methods **Contiguous** Allocation





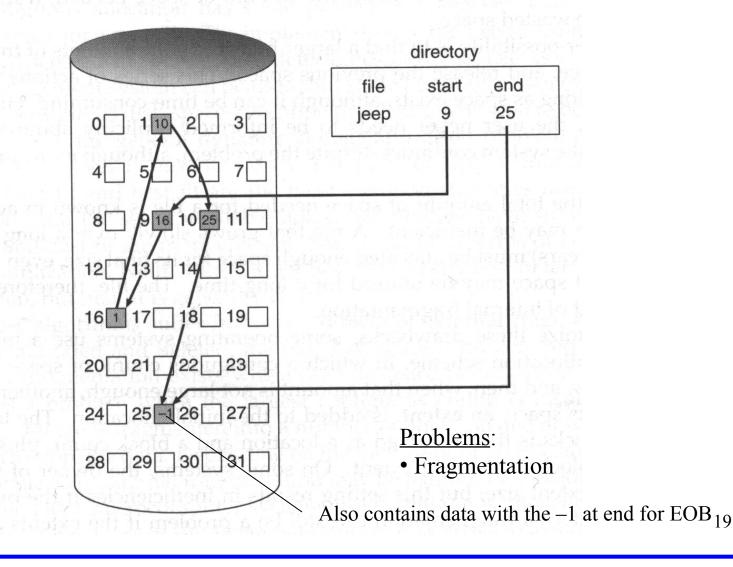
No EOB pointer needed (EOF used anyway)

#### Problems:

- Fragmentation
- Enough continuous space
- Multiple file handles (how much space to leave?)

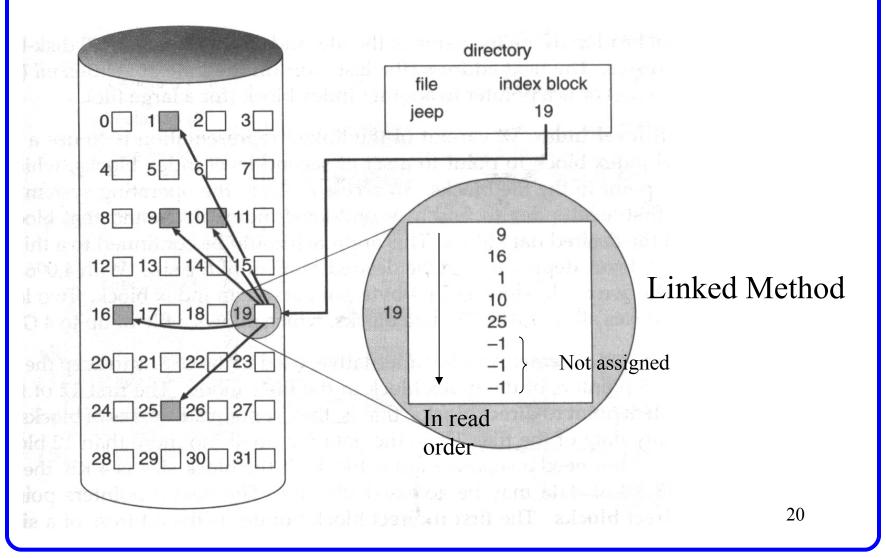


### File Allocation Methods Linked Allocation





### File Allocation Methods Indexed Allocation



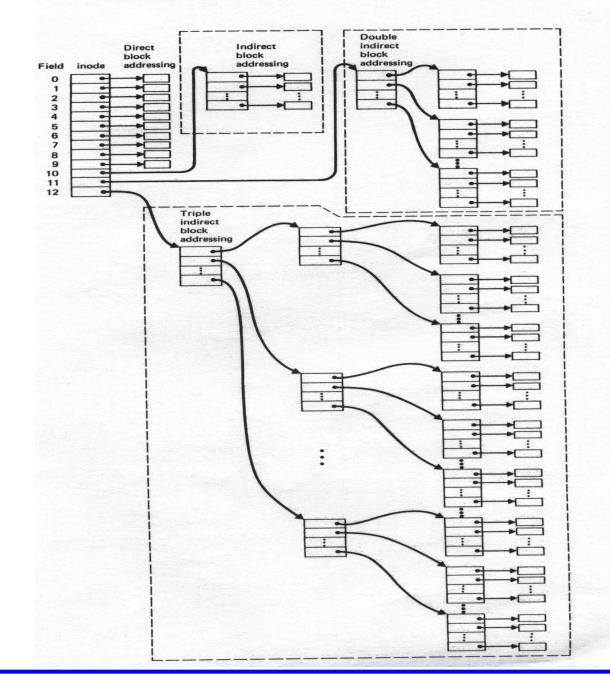
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- Multilevel Index
  - First index block is a pointer to a maximum of n simple Linked method blocks that point to files.
- Combined Index
  - For example the Unix File System does this:
    - 15 pointers to files
      - 12 directly to file data blocks
      - 1 points to a Linked block
      - 1 point to a Multilevel index block (2 levels)
      - 1 pointer to a Multilevel index block (3 levels)

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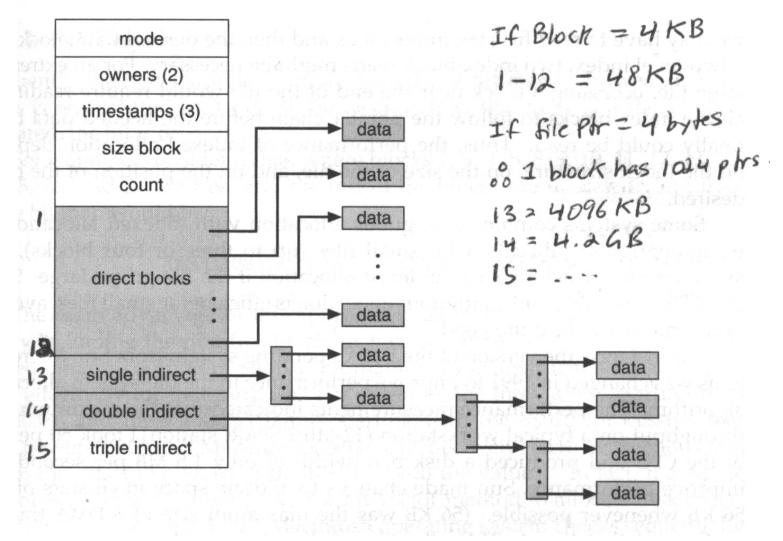




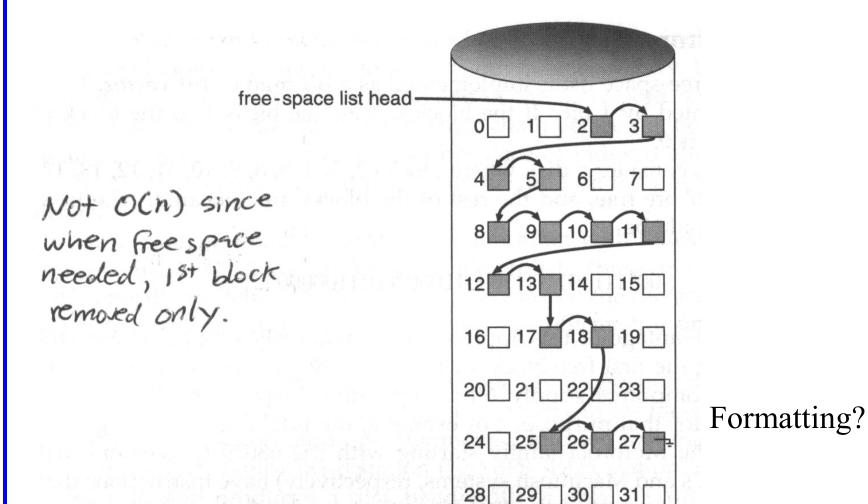
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## Example: Unix inode Method







#### Pointer Method

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# Managing Free Disk Space

The Bit Vector Method

- Stored on disk (but loaded in RAM to improve performance)
- 1 bit = 1 block (where 1 means free / 0 means used)

block No = (bits in word) X (No. of 0 words) + offset to 1 bit

000111110000001111100000 etc.

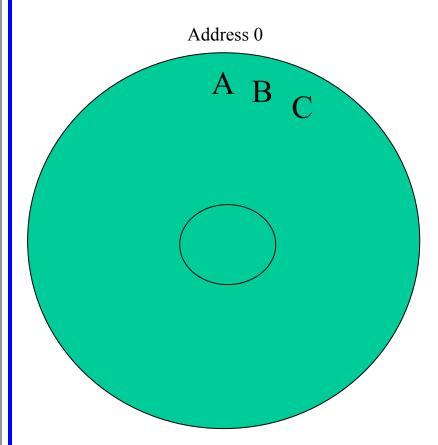
Eg/ The size of a Bit vector table? 1.3 GB HDD with 512 byte blocks = about 332KB (fast only if stored in RAM)(used by MAC)



## Part 3

### OS Management of File Systems

## Basic Disk Layout



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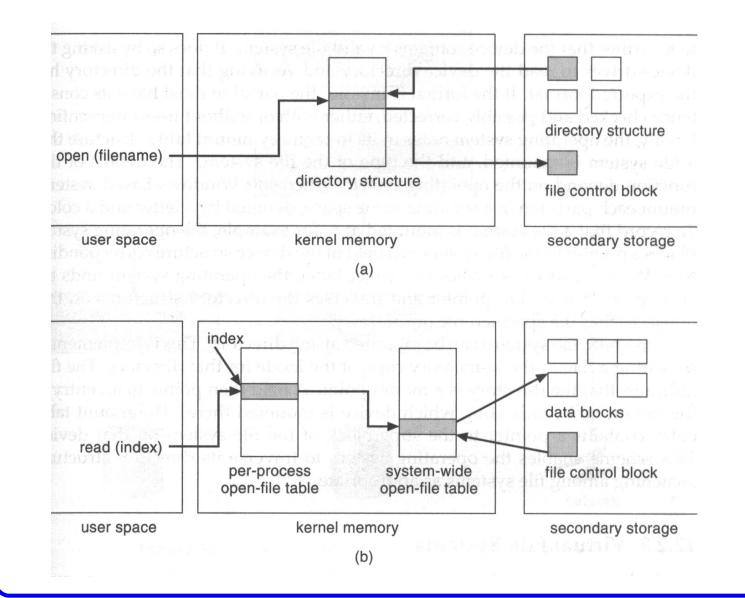
A = Boot Control Block B = Partition Control Block C = Directory Structure D = File Control Block (FCB) (either in C or external) ↑

Block = track/sector intersection or fixed byte size

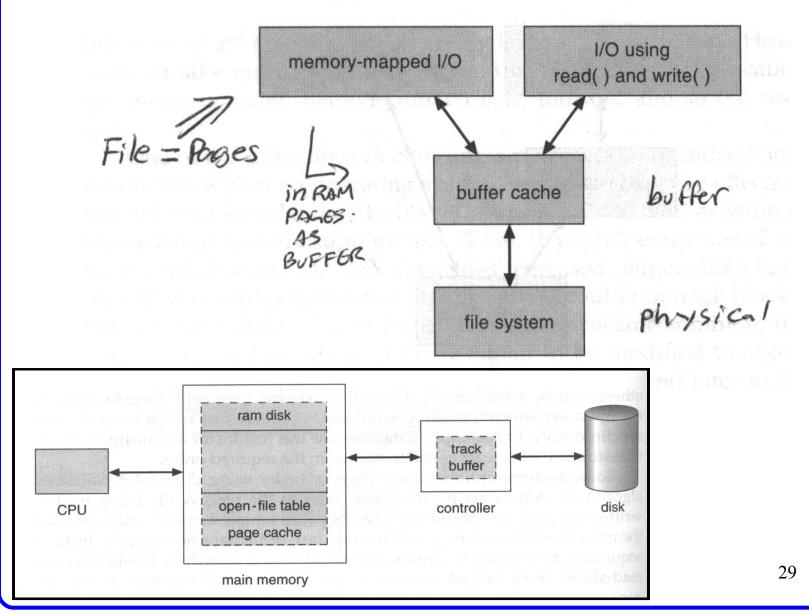
a.k.a FAT



# OS Management of Directories

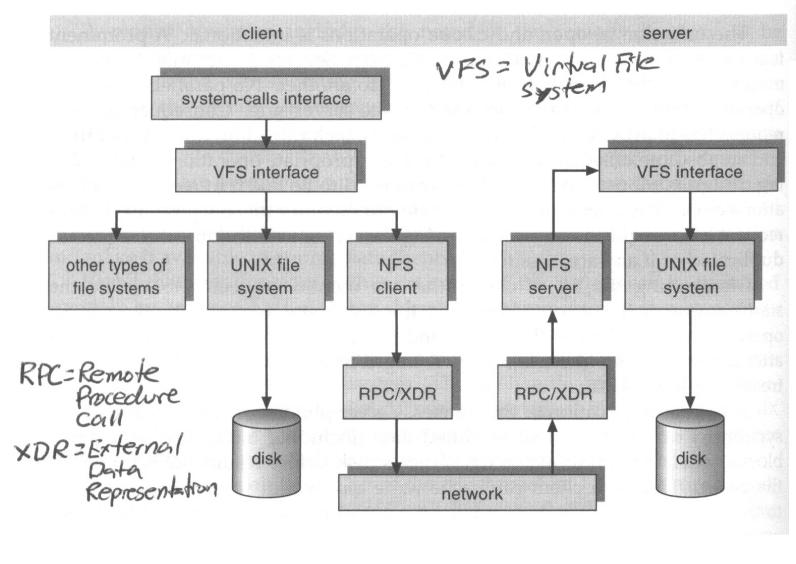


## Buffers & More Buffers...



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## Example NFS OS Architecture





## Question

• How could we implement social networking through directory structures?





## Part 4

#### At Home



# Things to try out

- Using the command line prompt mount a drive letter to an external network directory using Unix
- 2. Now do this using the window from the command-line interface and the windowed interface.
- 3. Web Resources:
  - http://www.beginningtoseethelight.org/fat16/
  - http://www.dewassoc.com/support/win98/fat.htm
  - http://home.no.net/tkos/info/fat.html